

Markus Schlegel, Active Group, 2025-05-07



Decoupled by Default

Funktionale Softwarearchitektur

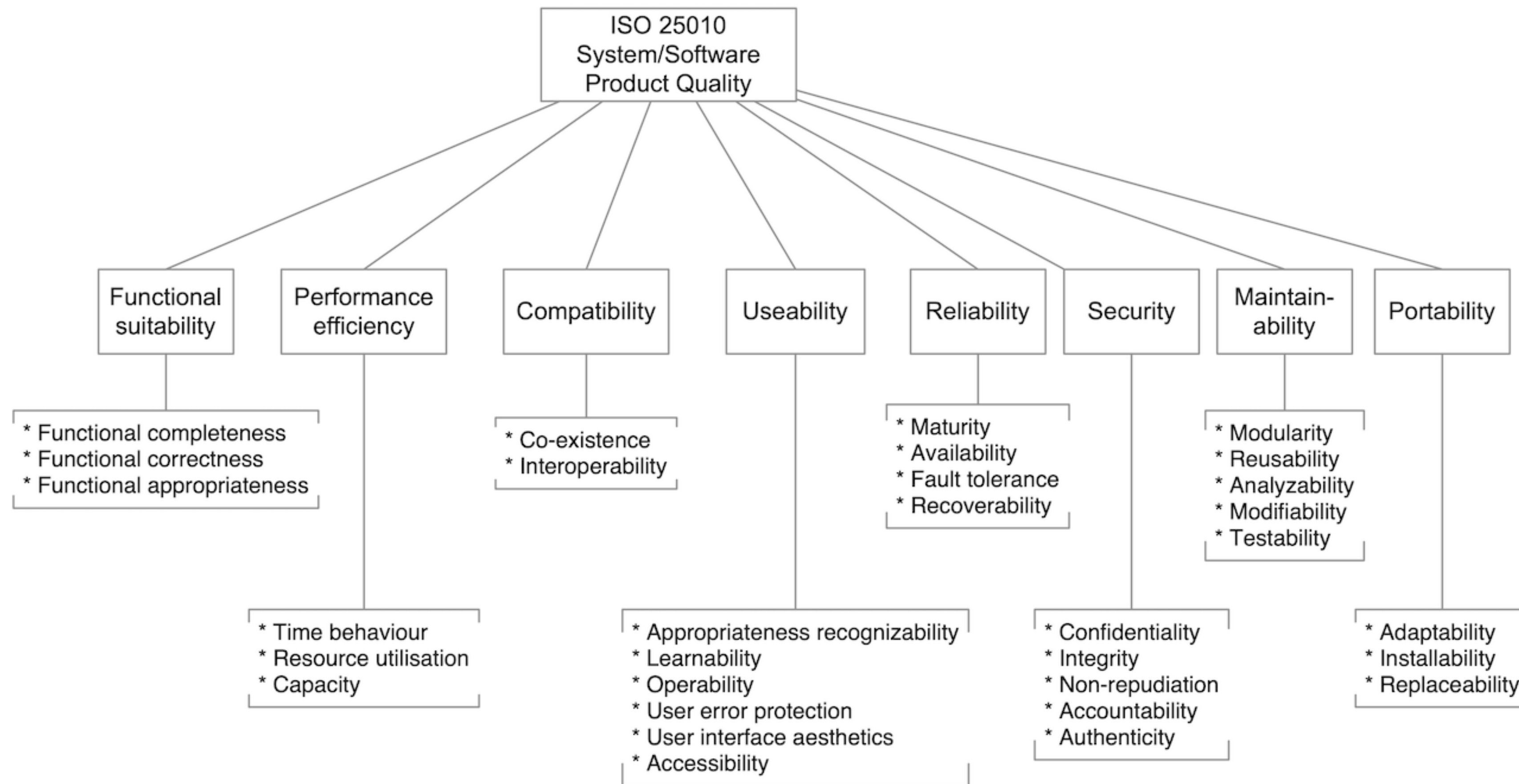
Funktionale Softwarearchitektur

"Functional programming in the large"

Functional Software Architecture refers to methods of construction and structure of large and long-lived software projects that are implemented in functional languages and released to real users, typically in industry.

<https://functional-architecture.org/>

Warum Wartbarkeit?



Ohne Wartbarkeit ist alles doof.

Funktionalität :

Zuverlässigkeit :

Sicherheit

Compatibility :

Portability

Krabbelkäfer

Lieblingsarchitekten

Performance

Zum Ausfüllen
und Verschenken
für deinen

riva

Decoupled (Coupling)

Dt. Kopplung

Kopplung ist das Maß der Abhängigkeiten zwischen zwei Modulen.

“Je mehr wir über Modul B wissen müssen, um Modul A zu verstehen, desto enger gebunden ist A an B.”

– Yourdon & Constantine, Structured Design

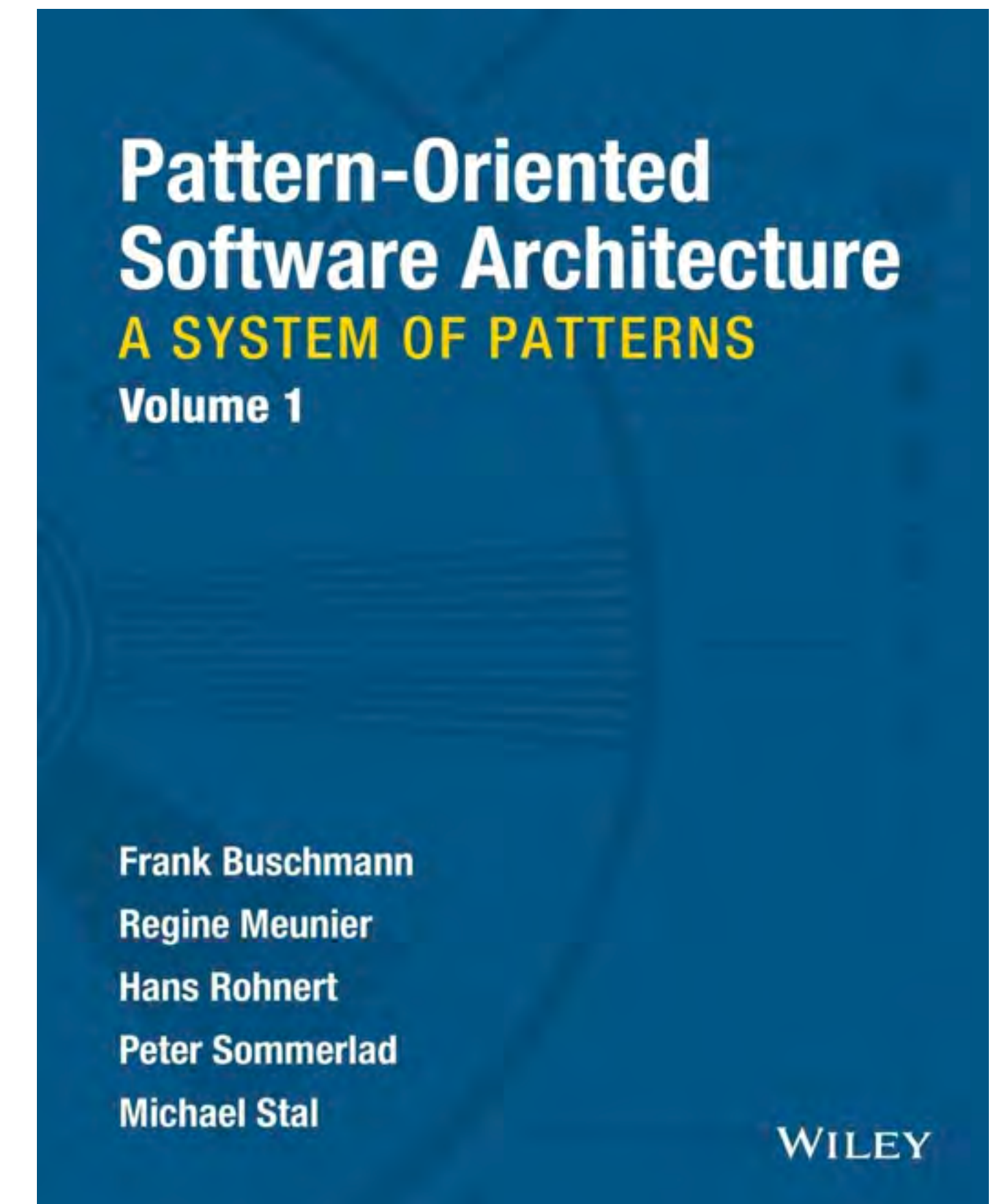
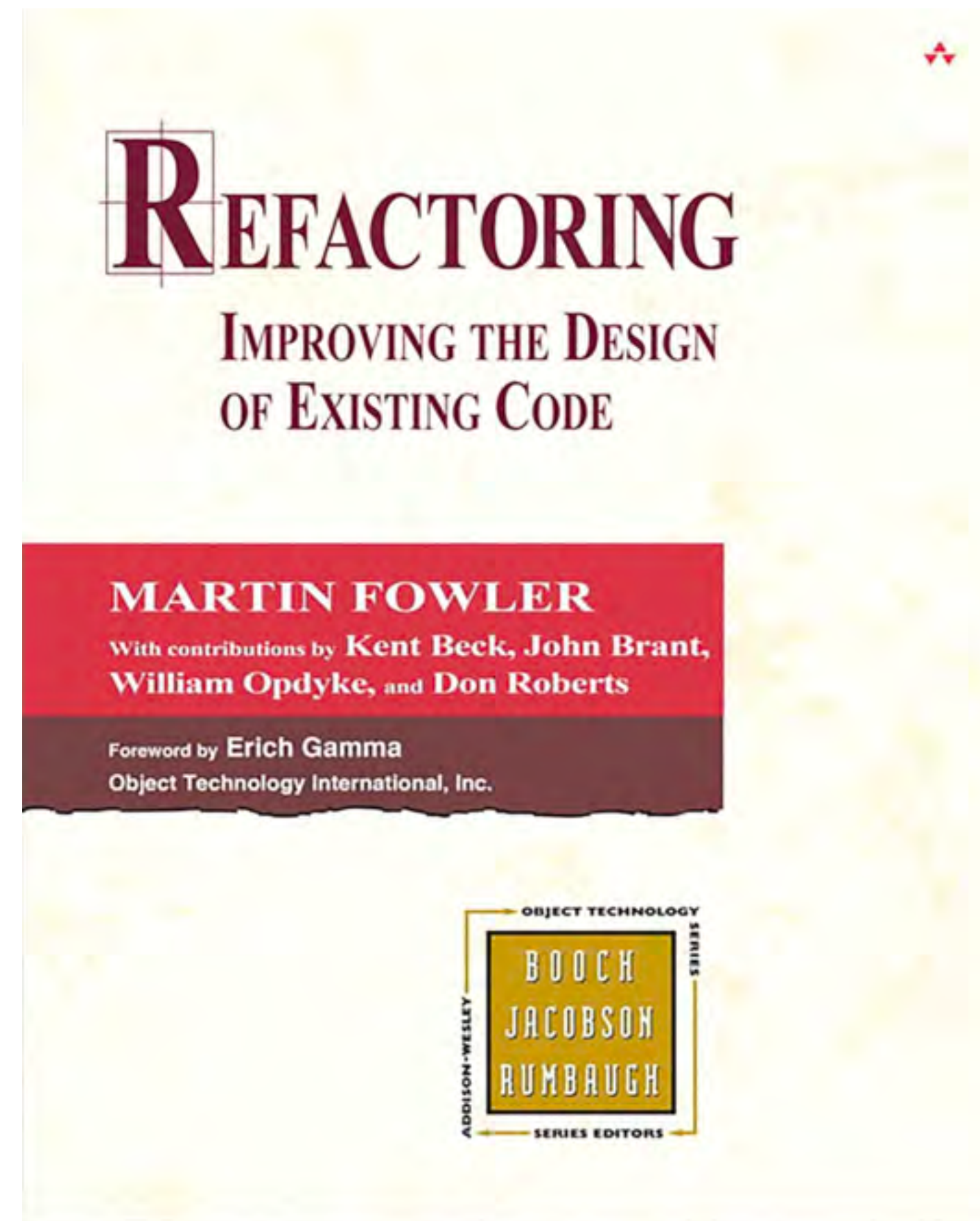
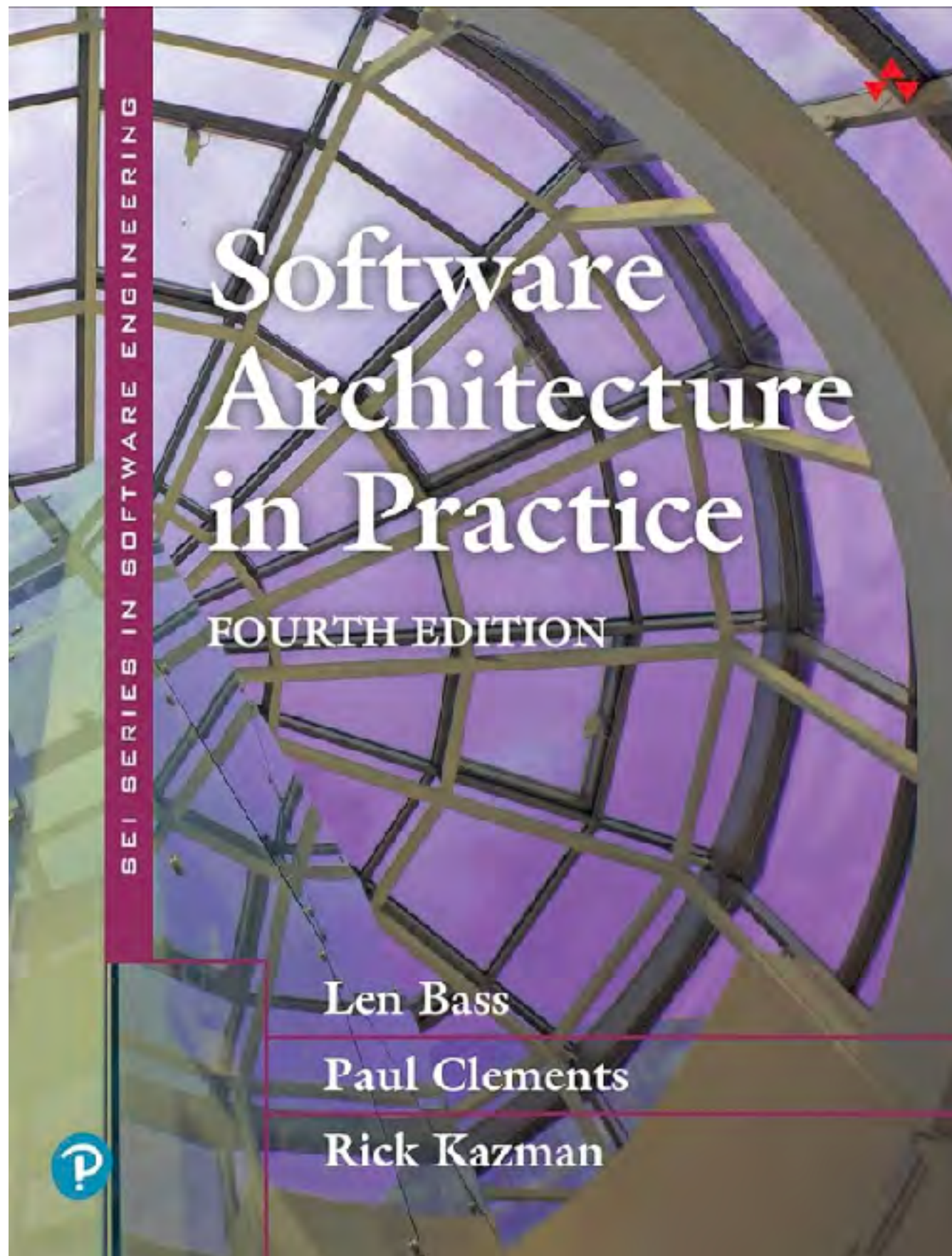


Decoupled (Coupling)

Dt. Kopplung

Kopplung ist damit das entscheidende Maß für die Wartbarkeit (Formbarkeit) eines Softwaresystems.

Decoupled **by Default**



Arten von Kopplung

- Common Coupling: geteilter Zustand
- External Coupling: Kopplung an die Plattform
- Content Coupling: Verletzung des Information-Hiding-Prinzips
- ...

Common Coupling

Programmieren mit veränderlichem Zustand

```
public class Person {  
    private String name;  
    private String address;  
    ...  
    public void setAddress(String newAddress) {  
        this.address = newAddress;  
    }  
}
```

```
public class A {  
    private Person myPerson;  
    ...  
    public void receivePerson(Person p) {  
        myPerson = p;  
    }  
}
```

```
public class B {  
    private Person myPerson;  
    ...  
    public void transferPerson(A other) {  
        // Transfer to other department  
        other.receivePerson(myPerson);  
    }  
  
    public void backToTheOffice() {  
        myPerson.setAddress("Seattle");  
    }  
}
```

Imperativ: Data

Funktional: Data - Mutation

Imperativ: Data

Funktional: Data - Mutation

Während *Daten* in der Umgangssprache Gegebenheiten, Tatsachen oder Ereignisse sind, sind *Daten* in der Fachsprache Zeichen, die eine Information darstellen.

Imperativ: Data

Funktional: Data - Mutation

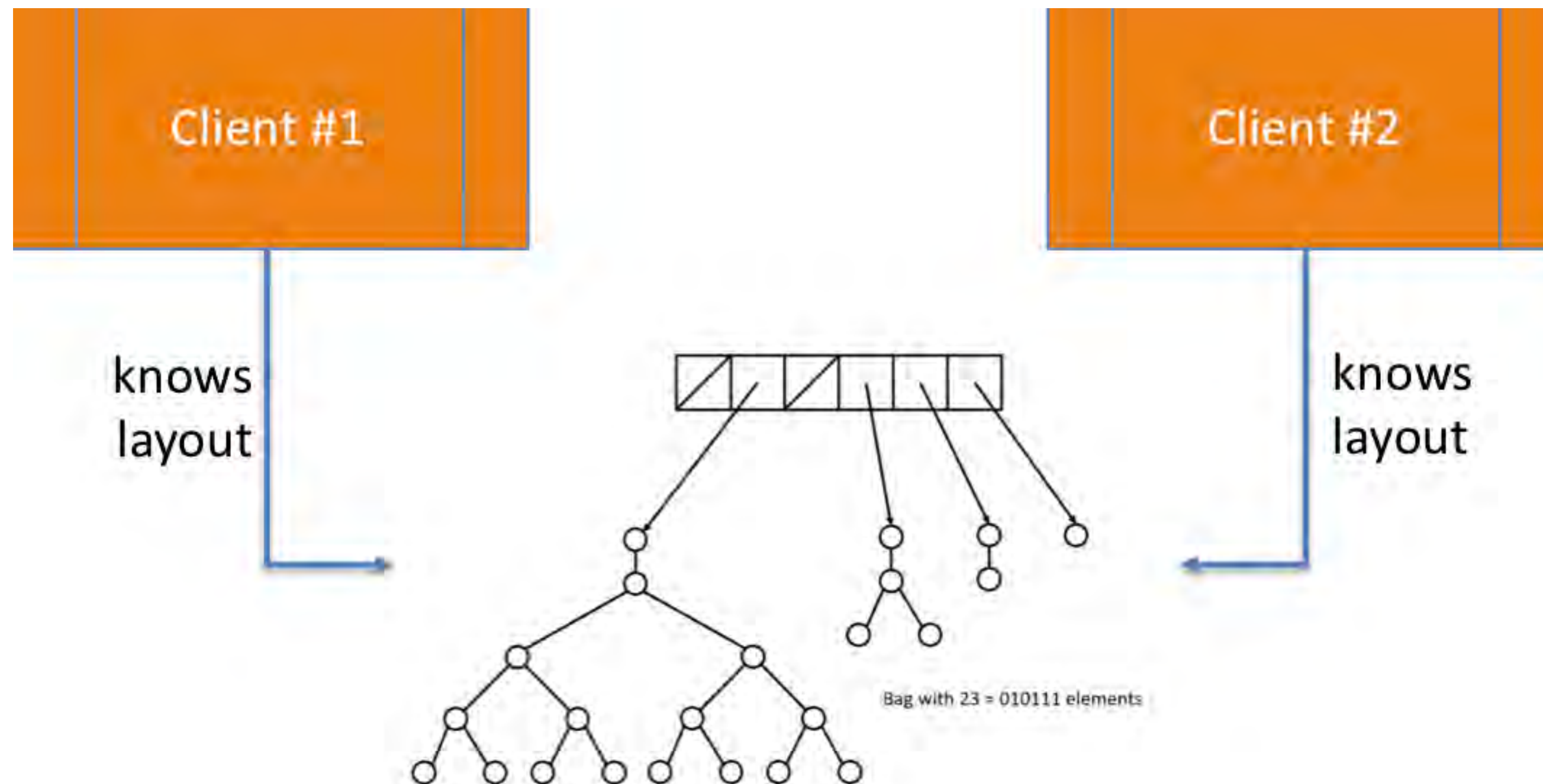
Während *Daten* in der Umgangssprache Gegebenheiten, Tatsachen oder Ereignisse sind, sind *Daten* in der Fachsprache Zeichen, die eine Information darstellen.

Imperativ: Data + Place

Funktional: Data

Content Coupling

Verletzung des Information-Hiding-Prinzips



Zeitreihen

```
interface TimeSeriesServiceLegacy {  
    public List<Pair<Time, Double>> getTimeSeriesData(Time from, Time to);  
}
```

```
var ts1 = listOf(new Pair<>(t1, 5.0),  
                 new Pair<>(t2, 6.5),  
                 new Pair<>(t3, 7.3));
```

```
var ts2 = listOf(new Pair<>(t2, 6.5),  
                 new Pair<>(t1, 5.0),  
                 new Pair<>(t3, 7.3));
```

```
var ts3 = listOf(new Pair<>(t1, 6.5),  
                 new Pair<>(t1, 6.5));
```

```
var ts4 = listOf(new Pair<>(t1, 6.5),  
                 new Pair<>(t1, 9.9));
```

Make Illegal States Unrepresentable

Effective ML video

AUG 21, 2010 | 1 MIN READ



By: [Yaron Minsky](#)

A while back I [mentioned](#) that I'd given a guest lecture at classes at Harvard and Northeastern, and that the Harvard class had been taped. I finally got and uploaded that video:



Make Illegal States Unrepresentable

Aber wie?

- **Smart Constructors**
- Parsing mit nominell frischen Typen (**“Parse, don’t validate”**, <https://lexi-lambda.github.io/blog/2019/11/05/parse-don-t-validate/>)
- Flache Produkttypen mit nullables zu **Summentypen** machen (<https://funktionale-programmierung.de/2024/11/25/sums-products.html>)
- Assoziative **Maps** nutzen
- **Funktionen** nutzen
- ...

```
type connection_state =  
| Connecting  
| Connected  
| Disconnected  
  
type connection_info = {  
  state: connection_state;  
  server: Inet_addr.t;  
  last_ping_time: Time.t option;  
  last_ping_id: int option;  
  session_id: string option;  
  when_initiated: Time.t option;  
  when_disconnected: Time.t option;  
}
```

```
type connecting = { when_initiated: Time.t; }  
type connected = { last_ping : (Time.t * int) option;  
                   session_id: string; }  
type disconnected = { when_disconnected: Time.t; }  
  
type connection_state =  
| Connecting of connecting  
| Connected of connected  
| Disconnected of disconnected  
  
type connection_info = {  
  state : connection_state;  
  server: Inet_addr.t;  
}
```

The whole is greater than the sum of the parts

- Borrowing some terminology from the functional world:
 - Tuples are sometimes called *product types*; the values of a tuple are drawn from the cartesian product of components types
 - A Point is an x *and* a y
 - Sealed classes (discriminated unions) are sometimes called *sum types*, because the values of a sealed class is the sum (union) of their permitted subtypes
 - A Shape is a Circle *or* a Rectangle
 - Product and sum types together are often called “algebraic data types” (ADTs)
- ADTs are a surprisingly powerful paradigm for describing data



JAVAONE'25

<https://www.youtube.com/watch?v=1dY57CDxR14>

```

public class TimeSeriesServiceAntiCorruption {
    private TimeSeriesServiceLegacy legacyService;

    public TimeSeriesServiceAntiCorruption(TimeSeriesServiceLegacy ls) {
        legacyService = ls;
    }

    private Map<Time, Double> parseList(List<Pair<Time, Double>> l, Map<Time, Double> acc) {
        if (l instanceof Nil) {
            return acc;
        } else if (l instanceof Cons<Pair<Time, Double>> cons) {
            return parseList(cons.cdr(),
                             acc.assoc(cons.car().left(),
                                         cons.car().right(),
                                         acc));
        } else {
            return null;
        }
    }

    public Map<Time, Double> getTimeSeriesData(Time from, Time to) {
        return parseList(legacyService.getTimeSeriesData(from, to), EmptyMap.getInstance());
    }
}

```

```

public class TimeSeriesServiceAntiCorruption {
    private TimeSeriesServiceLegacy legacyService;

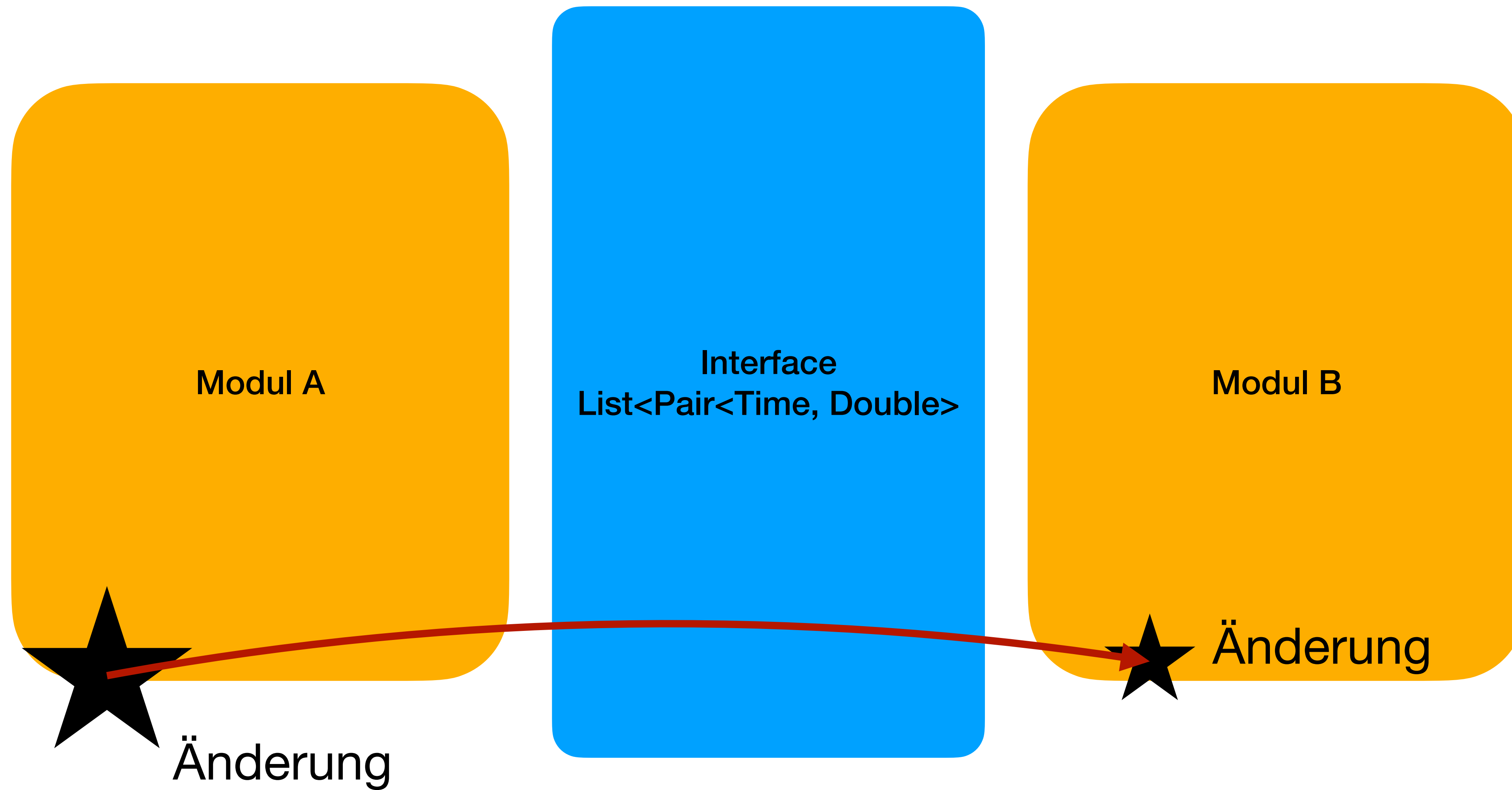
    public TimeSeriesServiceAntiCorruption(TimeSeriesServiceLegacy ls) {
        legacyService = ls;
    }

    private Map<Time, Double> parseList(List<Pair<Time, Double>> l, Map<Time, Double> acc) {
        if (l instanceof Nil) {
            return acc;
        } else if (l instanceof Cons<Pair<Time, Double>> cons) {
            return parseList(cons.cdr(),
                             acc.assoc(cons.car().left(),
                                         cons.car().right(),
                                         acc));
        } else {
            return null;
        }
    }

    public Map<Time, Double> getTimeSeriesData(Time from, Time to) {
        return parseList(legacyService.getTimeSeriesData(from, to), EmptyMap.getInstance());
    }
}

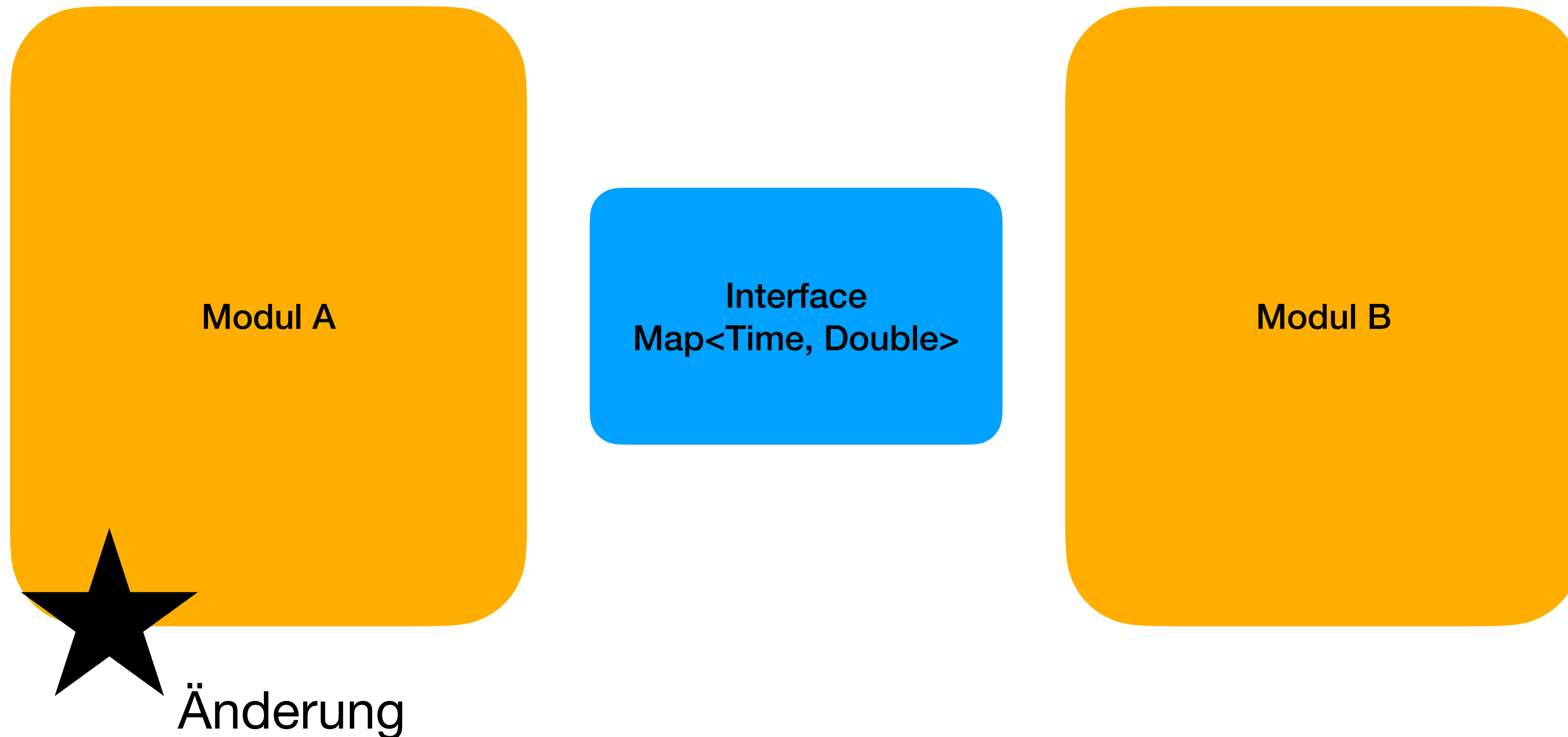
```

Zu offenes Interface



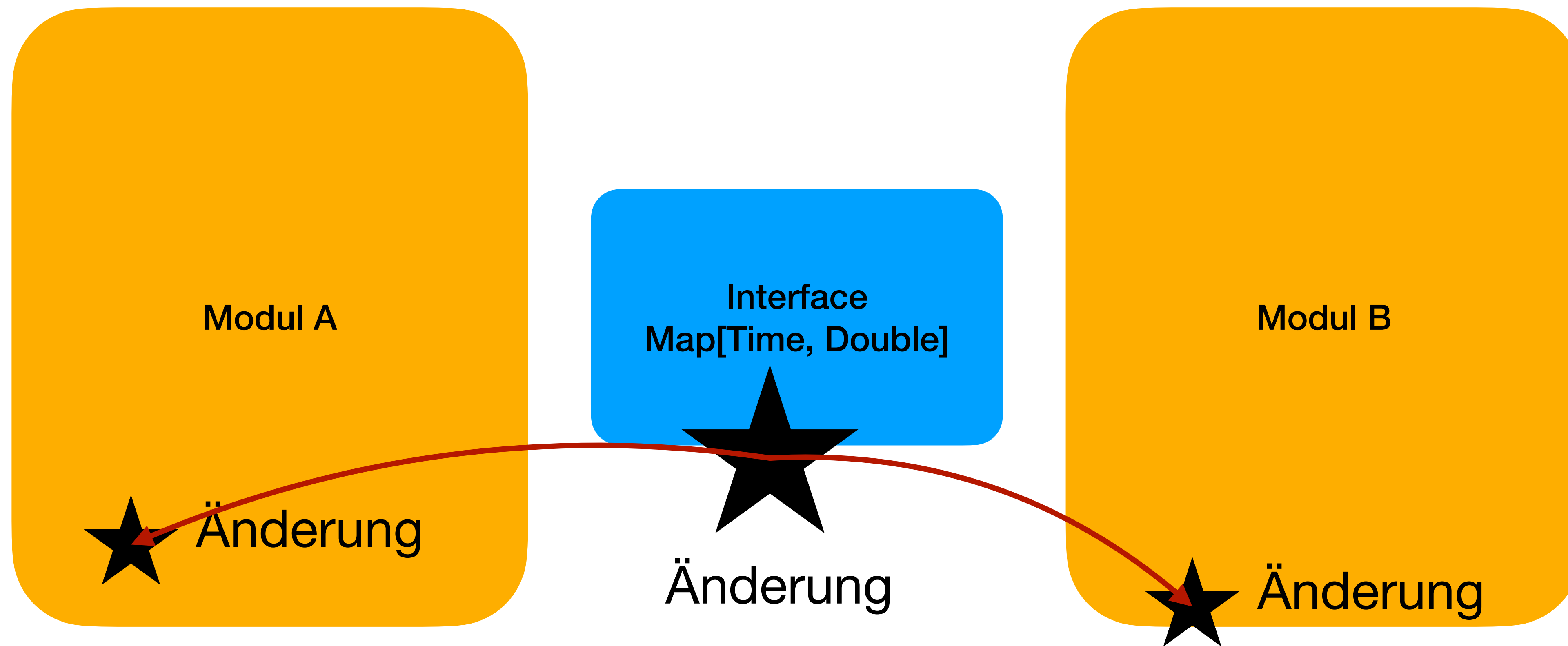
Zu beschränktes Interface

(Make Illegal States Unrepresentable)



Zu beschränktes Interface

(Make Illegal States Unrepresentable)



Make Legal States Representable

Denotational Design

LambdaPix technical report 2009-01, March 2009 (minor revisions January 28, 2016)

1

Denotational design with type class morphisms (extended version)

Conal Elliott

LambdaPix
conal@conal.net

Abstract

Type classes provide a mechanism for varied implementations of standard interfaces. Many of these interfaces are founded in mathematical tradition and so have regularity not only of *types* but also of *properties* (laws) that must hold. Types and properties give strong guidance to the library implementor, while leaving freedom as well. Some of this remaining freedom is in *how* the implementation works, and some is in *what* it accomplishes.

To give additional guidance to the *what*, without impinging on the *how*, this paper proposes a principle of *type class morphisms* (TCMs), which further refines the compositional style of denotational semantics. The TCM idea is simply that *the instance's meaning follows the meaning's instance*. This principle determines the

little. It is form without *essence*. Users of this abstraction care about what *functionality* these names have. They care what the names *mean*.

Which leads to the question: what do we mean by “mean”? What is the *essence* of a type, enlivening its form? One useful answer is given by denotational semantics. The meaning of a program data type is a type of mathematical object (e.g., numbers, sets, functions, tuples). The meaning of each operation is defined as a function from the meanings of its arguments to the meaning of its result. For instance, the meaning the type *Map* *k* *v* could be partial functions from *k* to *v*. In this model, *empty* is the completely undefined function, *insert* extends a partial function, and *lookup* is just function application, yielding $\perp$ where undefined. The meaning of the

Time Series = Map = Partial + Function ?

Time Series = ~~Map = Partial~~ + Function

```

interface TimeFunction<A> extends Function<Time, A> {}

record TFList<A>(List<Pair<Time, A>> list) implements TimeFunction<Maybe<A>> {
    public Maybe<A> apply(Time t) {
        if (list instanceof Nil) {
            return Maybe.nothing();
        } else if (list instanceof Cons<Pair<Time, A>> cons) {
            if (cons.car().left().equals(t)) {
                return Maybe.just(cons.car().right());
            } else {
                return new TFList<>(cons.cdr()).apply(t);
            }
        } else {
            return null;
        }
    }
}

```

```
public class TimeSeriesService2 {  
    private TimeSeriesServiceLegacy legacyService;  
  
    public TimeSeriesService2(TimeSeriesServiceLegacy ls) {  
        legacyService = ls;  
    }  
  
    public TimeFunction<Maybe<Double>> getTimeSeriesData(Time from, Time to) {  
        return new TFList<Double>(legacyService.getTimeSeriesData(from ,to));  
    }  
}
```

```
record TFFun<A>(Function<Time, A> f) implements TimeFunction<A> {  
    public A apply(Time t) {  
        return f.apply(t);  
    }  
}
```

```
var immerPi = new TFFun<Double>(t -> 3.1415);
```

Zusammenfassung

Decoupled by Default

- Wartbarkeit/Formbarkeit ist der Enabler für alle weiteren Softwarequalitäten
- Kopplung ist das wichtigste Maß für Wartbarkeit
- Funktionale Architekturen sind per Werkseinstellung wartbarer als klassische Architekturen
- Mit Denotational Design entstehen ganz beiläufig lose gekoppelte Systeme
- Mehr gibt's im iSAQB-Modul FUNAR: <https://active-group.de/schulung/FUNAR/>



Getypte funktionale Programmierung auf der Java-Plattform mit Scala

Scala – woher, wohin, warum?

Funktionale Programmierung ist besser

Mächtigere Modelle, weniger Kopplung, aussagekräftige Typen

Einführung in die funktionale Programmierung mit Scala

Von null auf Scala

Softwarearchitektur: Warum Kombinatormodelle die Domäne revolutionieren können

Kombiniere ... das Modell passt

Property-based Testing in Scala

Don't Write Tests: Write Properties!

Funktionale Programmierung in der Softwarearchitektur

Decoupled by Default

Mailand oder Neubauten, Hauptsache S(k)cala

Editorial